

# Peer-to-peer networks

Pioneers, self-organisation, small-world-phenomenons

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- 2 The first peer-to-peer networks
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# What is a peer-to-peer network

## Types of networks I



Figure: Client server based network, *Wikipedia.de*



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# What is a peer-to-peer network

## Types of networks II

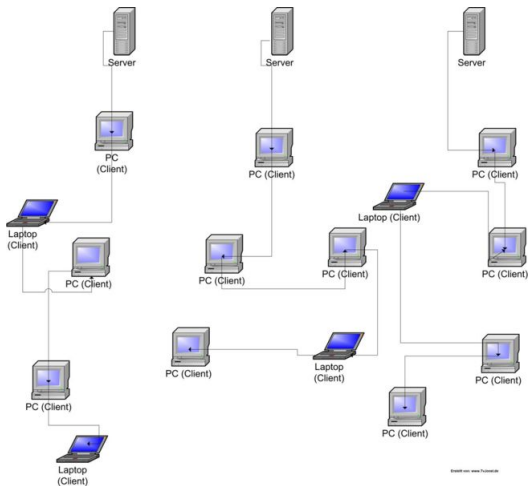


Figure: Peer-to-peer network, [Wikipedia.de](http://Wikipedia.de)



# What is a peer-to-peer network

## Definition

- A connection between several network participants, for exchanging data.
- Two participants directly communicate with each other.
- Nodes are called *peers*.
- All peers have equal rights.



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- Released by *Shawn Fanning* in 1999.
- Actually a very simple protocol.
- First massively popular peer-to-peer network (download of the year 2000).
- It allowed to share audio data in the *.mp3* format with thousands of other users.
- Because of offering copyright protected files, Napster was sued several times.
- Therefore, Fanning started a cooperation with Bertelsmann Ecommerce in the year 2000.
- Today Napster is based on a client-server architecture.



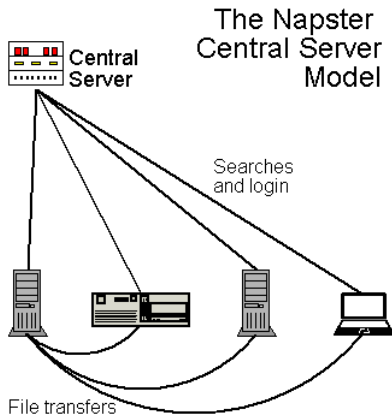


Figure: Napster infrastructure





- Napster is actually no peer-to-peer network (a server is necessary!).
- Peer tells server which media it offers and its address.
- Server keeps track of all peers and their media files.
- Procedure:
  - Peer contacts server with a request for a file.
  - Server responds with a list of addresses containing peers sharing the requested media.
  - The requesting peer contacts a peer from the list it receives and downloads the file directly.



- Assets:
  - Simple design concept.
  - Good solution for distributing files over several nodes.
- Drawbacks:
  - Server is a single point of failure.
  - Protocol is not scalable.



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- Developed by *Justin Frankel* and *Tom Pepper* in early 2000.
- Gnutella protocol is open-source.
- Therefore, many different clients arose within the next years.
- Most famous clients: *LimeWire* and *Morpheus*.



- Gnutella is a real peer-to-peer network (there exists no server).
- The protocol has an important parameter (*TTL*) to control the network structure.
- First problem is that the connection setup for a peer is complicated.
- Which peer should be called on connection setup, avoiding that the connected peer becomes a server?



### ■ Solution: *Bootstrapping*

- 1 Initially a peer has a preexisting list with peers.
- 2 On startup, peer contacts all other peers from list until an active peer answers.
- 3 The contacted peer sends the message further and so on, until peer *TTL* away is reached.
- 4 Every peer on this path sends a response back to the starting peer.
- 5 Starting peer updates his neighborhood list for the next start.



- Use of five basic messages while operating (used on top of *TCP*):
  - Ping** Used for finding other nodes.
  - Pong** Response to the ping-message.
  - Query** For querying the network for files.
  - QueryHit** Answer to a query, when the node shares the queried file. Contains the IP and port address of the sender.
  - Push** Technical message for sharing behind firewalls.



- After sending the Ping-message on startup, the peer receives Pong-messages.
- The peer chooses randomly from the answering peers  $k$  as its neighbors.
- A typical value for  $k$  is 5.
- The Query-message is send to all neighbors and so do they, until peer  $TTL$  away is reached.
- If one of this peers answers with a QuerHit-message a direct connection is established and this is added as neighbor.





- Assets:
  - No single point of failure (no server).
  - Scalable and stable structure.
- Drawbacks:
  - Peer can only reach the peers *TTL* away.
  - Therefore, rare files might not be found in large networks.



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*"...the ability of systems comprising many units and subject to constraints, to **organize themselves** in various spatial, temporal or spatiotemporal activities. These emerging properties are pertinent to the system as a whole and cannot be seen in units which comprise the system..."*

- A. Babloyantz



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# Pareto distributions

## Short introduction

- Several measurements have analyzed the average number of neighbors of a peer in a Gnutella network.
- As a result, the average number of peers with  $d$  neighbors can be approximated by  $\frac{C}{d^k}$ , with  $C$  and  $k$  as network specific constants.
- Such a relation is called a *power law*...
- ... and the according probability distribution is called a *pareto distribution*.
- Pareto distributions are an often observed phenomenon in social processes.
- **Remember:** This is not the result of an algorithm, it actually is the result of a social phenomenon.



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# Pareto distributions

## The diameter of Gnutella

- Only asking the next *TTL* neighbors, we must consider the diameter of the network.
- Actually five measurements in the year 2000 have shown:

Diameter of Gnutella  $\approx 8 \dots 12$  peers

- This also can be explained by a social phenomenon.



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# Small world phenomens

## Description I

- Idea: Every member in a social networks is connected with every other member only over a small chain of members.
- Phenomen was first described by *Stanley Milgram* who made the following experiment:
  - Some letters were given to 60 people in Omaha, Nebraska addressed to a special destination in Sharon, Massachusetts.
  - The letter could only be forwarded to known persons.
- Result:
  - Most of the letters actually arrived.
  - The average number of station passed was only 5.5 .





- Networks with a small diameter are therefore called small-world networks.
- To understand how they emerge, we look at three modeling approaches:
  - Watts und Strogatz's approach
  - Kleinberg's approach
  - Barabasi und Albert's approach



# Small world phenomens

## Watts und Strogatz's approach I

- Starting with a ring network with  $n$  nodes.
- Every node is connected to the next  $k/2$  neighbors to the left and the right.
- Network consists of cliques and has a relatively large diameter.
- Now we replace every edge with the probability  $p \in [0, 1]$  by a random edge, leading to a random node.
- Result (for small  $p$ ):
  - Most of the cliques persist.
  - The diameter of the network decreases significantly.



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# Small world phenomena

## Watts und Strogatz's approach II

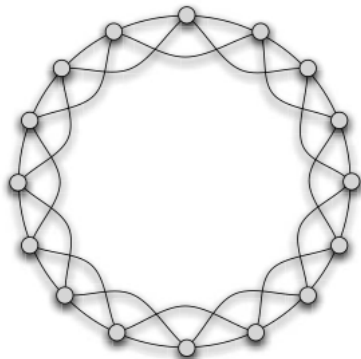


Figure: Starting with  $k = 4$



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# Small world phenomens

## Watts und Strogatz's approach III

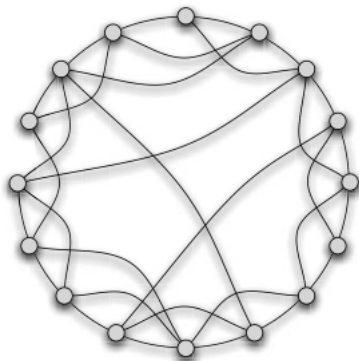


Figure: After alternating some edges



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# Small world phenomens

## Kleinberg's approach I

- Network is a grid network.
- Every node is connected to its direct neighbor.
- Special distant edges are added according to a special probability distribution.
- Kleinberg could prove that the diameter of the network is within  $\mathcal{O}(\log^2 n)$ .



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# Small world phenomens

## Kleinberg's approach II

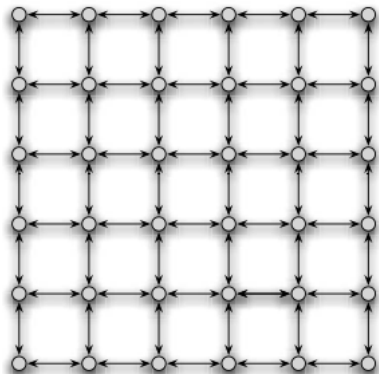


Figure: Starting position



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# Small world phenomena

## Kleinberg's approach III

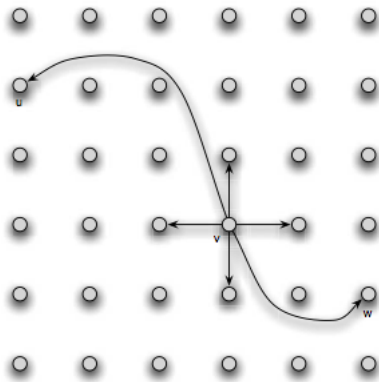


Figure: Adding distant edges



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# Small world phenomens

Barabasi und Albert's approach I

- Network is a small arbitrary graph.
- New nodes are added with  $m$  (as a constant) edges.
- Edges of the new nodes lead to old nodes according to a probability distribution, in which nodes, which already have many edges, are more likely to get a new one (rich gets richer).
- They could prove that the diameter of such a network is within  $\mathcal{O}(\log n)$ .



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# Small world phenomens

## Barabasi und Albert's approach II

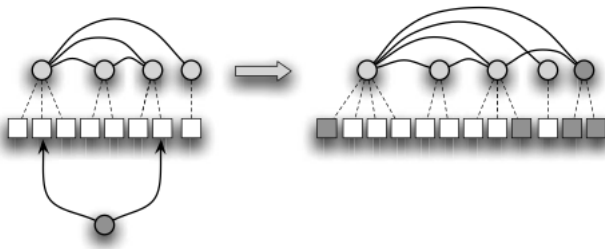


Figure: Adding a node



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# Small world phenomens

## Gnutella and small-world networks

In the year 2000 several measurements took place, which compared these three approaches with the actual Gnutella network, regarding the average distance between two peers. As a result we can say that:

**Barabasi und Albert** Biggest correspondence with Gnutella. This is justified by the similar methods used in this approach and the real Gnutella protocol.

**Watts Strogatz** Only moderate correspondence with Gnutella.

**Kleinberg** Only few correspondence with Gnutella.



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- Napster and Gnutella aren't in use any more.
- But the basic concepts of Napster and Gnutella are still in use today.
- Gnutella was detached by *Kademila*.
- Kademila introduced the concept of distributed hash tables.
- Kademila is still popular today. Famous clients: *eDonkey*, *BitTorrent* and *Azureus*.

